

On hereditarily finitely based varieties of semigroups

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Dedicated to Prof. H. Grell on his 70th birthday

In the investigation of the lattice of semigroup varieties one of the basic questions is the description of those varieties which have the property that all their subvarieties are finitely based. In particular, this means finding finite systems of identities which cannot be extended to infinite independent systems. P. PERKINS [4] has shown that commutativity has this property and the author of the present paper has extended this result to a large class of permutative identities [5] (as L. SHEVRIN informed me recently the same generalization had been obtained by A. AĪZENŠTAT).*) However, E. LYAPIN has disappointed those who had hoped to go far by this way: in [3] he has shown that "most" balanced identities can be included in infinite independent systems. Here we continue the work in this direction: we are going to show that semigroup identities which define hereditarily finitely based varieties belong to a few exceptional types.

§ 1. The main results

Consider the free semigroup F and the free monoid F^0 over a countably infinite alphabet $X = \{x_i | i = 1, 2, \dots\}$. The elements of X will be occasionally denoted also by x, y, z, y_i, z_i . A word is an element of F^0 , its identity element being the empty word \emptyset . If $u, v \in F^0$ and there exist two further words u', u'' such that $v = u'uu''$ then u is a part of v . If $u' = \emptyset$ or $u'' = \emptyset$ then u is a beginning part or an end part, respectively.

Define the quasi-order \triangleleft on F by

Definition 1. $u \triangleleft v$ ($u, v \in F$) iff there exists an endomorphism $\varphi: F \rightarrow F$ such that $u\varphi$ is a part of v .

If we denote the length of a word w by $l(w)$ then $u \triangleleft v$ obviously implies $l(u) \leq l(v)$. — Now we extend the relation \triangleleft to subsets of F .

*) Remark at sheet-proof: As a matter of fact, this result has been obtained already by PUTCHA and YAQUB (*Semigroup Forum*, 3 (1971), 68—73) although they have not formulated it explicitly.

Definition 1'. $U \triangleleft V$ ($U, V \subseteq F$) iff $u \triangleleft v$ for at least one pair $(u, v) \in U \times V$. In this case we say that V depends on U ; in the opposite case it is *independent from* U .

Note that for subsets the relation \triangleleft is not a quasi-order. — If $U = \{u\}$ we shall write also $u \triangleleft V$.

Definition 2. The set $V \subseteq F$ is *dependent* (in itself) if $v \triangleleft V \setminus \{v\}$ for some $v \in V$, and *independent* in the opposite case.

Definition 3. A finite subset $U \subset F$ is an *essentially finite set* (EFS) if $U \triangleleft V$ holds for every infinite independent $V \subseteq F$.

The significance of these notions for our purpose is established by the following propositions.

Proposition 1. *If the (finite or infinite) system of non-trivial semigroup identities*

$$(\tau) \quad v_{2k-1} = v_{2k} \quad (k = 1, 2, \dots)$$

is such that $V = \{v_i | i = 1, 2, \dots\}$ is an independent set of words then (τ) is an independent system of identities.

Proof. Denote by K_i the characteristic ideal of F generated by all v_j 's, $j \neq i$, i.e. the ideal generated by $\{v_j \varphi | j \neq i, \varphi \in \text{End}(F)\}$. It is easy to see that for every homomorphism $\chi: F \rightarrow F/K_i$ we have $v_j \chi = 0$ for $j \neq i$ so that all identities of (τ) but the one containing v_i hold in F/K_i . On the other hand, the independence of V means that $v_i \notin K_i$ so that the element \bar{v}_i of F/K_i corresponding to v_i under the natural homomorphism is not 0. Thus, no identity in (τ) follows from the others.

Proposition 2. *Let (τ) be as in Proposition 1 but infinite and*

$$(\sigma) \quad u_s = u'_s \quad (s = 1, \dots, m)$$

arbitrary. If $V = \{v_i | i = 1, 2, \dots\}$ is independent from $U = \{u_i, u'_i | 1 \leq i \leq m\}$ then $(\sigma) \cup (\tau)$ has no finite basis.

Indeed, as above, none of the identities in (τ) follows from the rest of $(\sigma) \cup (\tau)$. However, if a finite basis existed, it could be chosen as a subsystem of $(\sigma) \cup (\tau)$.

Corollary. *If \mathfrak{S} is a hereditarily finitely based variety of semigroups and (σ) is a basis of \mathfrak{S} then U is an EFS.*

Indeed, in the opposite case an infinite independent system V would exist such that U non $\triangleleft V$ and the subvariety of \mathfrak{S} defined by $(\sigma) \cup (\tau)$ would have no finite basis in virtue of Proposition 2.

By this Corollary, if we succeeded in determining all essentially finite subsets of F we could attain a considerable restriction of the scope of varieties which may be hereditarily finitely based. In this paper we determine all EFS's with 1 or 2 elements.

It holds obviously:

Proposition 3. *If $U \subseteq U'$ and U is an EFS then so is U' .*

More generally,

Proposition 4. *If U is an EFS and U' is a finite set of words such that $U' \triangleleft \{u\}$ for every $u \in U$ then U' is an EFS, too.*

Indeed, every set depending on U depends on U' ; so do all infinite independent sets.

In what follows we need some further definitions.

Definition 4. Two words u, v are *relatively prime* if no letter occurs in both of them.

Definition 5. u is a *closed part* of v if $v = u'uu''$ and u is relatively prime to u', u'' . If $u \in X$ it is said to be a *closed letter*.

Definition 6. The decomposition $v = v_1 \dots v_k$ is *closed* if every pair v_i, v_j ($i \neq j$) is relatively prime. (Notation: $v \doteq v_1 \dots v_k$.)

Definition 7. A *type* T is a subset of F consisting of all automorphic images of any of its own elements.

In other words T consists of all elements of F which differ from each other only by the notation of letters. If $u \in T$ we shall write also $T = T(u)$.

Definition 8. A word is *simple* if all its letters are closed (i.e. if they are all distinct). Denote the set of nonempty simple words by X^* .

Definition 9. Let T be a type. The word v is *T -simple* if $v \doteq \prod_{i=1}^n v_i$ where $v_i \in T \cup X^*$. Denote the set of all nonempty T -simple words by T^* .

In what follows, for $U \subseteq F$ we shall denote $U \cup \{\emptyset\}$ by U^0 . Put furthermore $T_0 = \{v | v \doteq \prod_{i=1}^k v_i, v_i \in T\}$ and $T_1 = T^* \setminus T_0$.

Now we formulate the main results of our paper.

Theorem 1. *The one-element set $\{u\}$ is an EFS iff $u \in T(xyxx) \cup X^*$.*

Theorem 2. *The two-element set U is an EFS iff its elements are of type T and T' , respectively, where one of the following cases holds (v, v', w, w' always denote closed parts of the corresponding words, $w, w' \in X^{*0}$):*

- | | | | |
|-------------------|------------------|---------------------------------|---|
| (a) | $T = T(w),$ | T' arbitrary, | $w \neq \emptyset;$ |
| (b) | $T = T(xyxx),$ | T' arbitrary; | |
| (c) | $T = T(xwx),$ | $T' = T(x_1y_1x_1w'x_2y_2x_2),$ | $w, w' \neq \emptyset;$ |
| (d) | $T = T(xyzx),$ | $T' = T(vwv'),$ | $w \neq \emptyset, v, v' \in T(x^2)^0;$ |
| (d') | $T = T(xyzx),$ | $T' = T(x^2);$ | |
| (e ₁) | $T = T(xyzx),$ | $T' = T(x^2wyzy);$ | |
| (e ₂) | $T = T(xyzx),$ | $T' = T(yzywx^2);$ | |
| (f) | $T = T(wxyxw'),$ | $T' = T(v),$ | $v \in T^*(x^2);$ |

(g ₁)	$T = T(xywxw),$	$T' = T(xvyyx),$	$v \in T^*(x^2);$
(g ₂)	$T = T(wxyyx),$	$T' = T(xyyvx),$	$v \in T^*(x^2);$
(h ₁)	$T = T(xywxz),$	$T' = T(xvx),$	$v \in T_1(x^2);$
(h ₂)	$T = T(zxyyx),$	$T' = T(xvx),$	$v \in T_1(x^2).$

§ 2. Only if

In order to prove the “only if” parts we are going to list eight infinite independent word sets (Proposition 5); we shall see that all one- and two-element sets which the infinite subsets of the sets V_1, \dots, V_6'' depend on are those mentioned in the theorems. — For not to be obliged to prove the independence of each V_i separately we shall use the following

Lemma. Let u, u' and w be words having the following properties:

- i) *their first letters coincide with the last ones;*
- ii) *$l(w) > 1$;*
- iii) *if $l(u) > 1$ ($l(u') > 1$) then $u \text{ non } \triangleleft w$ ($u' \text{ non } \triangleleft w$);*
- iv) *if $l(u) = 1$ ($l(u') = 1$) then the letter $u(u')$ occurs in $u'(u)$.*

Then the set

$$V = \left\{ v_n | v_n = u \left(\prod_{i=1}^n w_i \right) u', w_i \in T(w), w_i \text{ closed} \right\}$$

is independent.

Sketch of the proof. Suppose $v_n \varphi$ is a part of v_m for $n < m$. First one shows that $u\varphi = u$ and $u'\varphi = u'$ (this follows from i, iii, iv and the fact that w_i is closed); hence $v_n \varphi = v_m$ and $l(w_k \varphi) > l(w)$ for some $k \leq n$. However, this is impossible by i, ii and the same fact as before.

Now we obtain immediately

Proposition 5. Put $a_n = \prod_{i=1}^n x_{2i-1} x_{2i} x_{2i-1}$, $b_n = \prod_{i=1}^n x_i^2$. Then the sets

$$V_1 = \{x_1 \dots x_{n-1} x_n x_{n-1} \dots x_1 x_n | n=3, 4, \dots\},$$

$$V_2 = \{y_1 z_1 z_2 y_1 a_n y_2 z_3 z_4 y_2 | n=1, 2, \dots\},$$

$$V_3 = \{y a_n y | n=1, 2, \dots\},$$

$$V_4 = \{y_1^2 a_n y_2^2 | n=1, 2, \dots\},$$

$$V_5 = \{y_1 z_1 y_1 b_n y_2 z_2 y_2 | n=1, 2, \dots\},$$

$$V_6 = \{y b_n y | n=1, 2, \dots\},$$

$$V_6' = \{y z y b_n z | n=1, 2, \dots\},$$

$$V_6'' = \{z b_n y z y | n=1, 2, \dots\}$$

are independent.

Proof. For V_1 this is known [1], for the rest it follows from the Lemma.

Before reverting to our Theorems, we state some simple propositions which we shall use later without referring to them.

Proposition 6. *If $u \in X^*$ then $u \triangleleft v$ iff $l(u) \equiv l(v)$.*

Proposition 7. *$xyx \text{ non } \triangleleft v$ iff $v \in T^*(x^2)$.*

Proposition 8. *$v \in T^*(xyx)$ iff $x^2 \text{ non } \triangleleft v$, $xyzx \text{ non } \triangleleft v$.*

Proposition 9. *If $u\varphi = v_1v_2v_3$, v_2 closed, then $u = u_1u_2u_3$, u_2 closed, $u_i\varphi = v_i$ ($i=1, 2, 3$).*

Now we can prove the "only if" parts of both theorems.

Necessity (Theorem 1). If $\{u\}$ is essentially finite and u is not simple then $u \triangleleft V_1$ implies $x^2 \text{ non } \triangleleft u$. Hence, by $u \triangleleft V_6$, u must be of the form xwx , w simple. However $u \triangleleft V_4$ implies $l(w) = 1$.

Necessity. (Theorem 2). Suppose $U = \{u_1, u_2\}$ is a minimal EFS (i.e. $u_1, u_2 \notin T(xyxx) \cup X^*$). Put first $x^2 \triangleleft u_2$. Then $U \triangleleft V_1$ implies $x^2 \text{ non } \triangleleft u_1$. By the same reason, no letter can occur in u_1 more than twice. Moreover, $u_1 \triangleleft V_1$, $u_1 \triangleleft V_2$ imply that at most (and then, by minimality of U , exactly) one letter may occur twice. Indeed, if x_i and x_j both occur twice in u_1 and say, the first occurrence of x_i precedes that of x_j then either this latter precedes the second occurrence of x_i ($u_1 = \dots x_i \dots x_j \dots \dots x_i \dots$) and $u_1 \text{ non } \triangleleft V_2$ or else (i.e. if $u_1 = \dots x_i \dots x_i \dots x_j \dots x_j \dots$) $u_1 \text{ non } \triangleleft V_1$. Thus, $u_1 = wx_iw'x_jw''$ (w, w', w'' simple and closed). Furthermore $u_1 \triangleleft V_2$ entails $l(w') \equiv 2$ and, since $u_1 \triangleleft V_3$, we have $w = w'' = \emptyset$ if $l(w') > 1$. Hence $u_1 = x_i x_j x_k x_i$ or $u_1 = wx_i x_j x_i w''$.

In the first case $u_1 \text{ non } \triangleleft V_4$, $u_1 \text{ non } \triangleleft V_5$ and so $u_2 \triangleleft V_4$ which implies $u_2 \in T(x^2)$ — case (d') — or $u_2 = v_1 q v_2$ with $v_1, v_2 \in T(x^2) \cup T(xyxx) \cup X$, $x^2 \text{ non } \triangleleft q$, $q \neq \emptyset$ if $v_1, v_2 \in T(x^2)$, and $u_2 \triangleleft V_5$ which implies $xyxx \text{ non } \triangleleft q$, $q \neq \emptyset$ if $v_1, v_2 \in T(xyxx)$. Now $x^2 \text{ non } \triangleleft q$, $xyxx \text{ non } \triangleleft q$ give $q \in X^{*0}$ and we get either one of the cases (d), (d'), (e₁), (e₂) or a subcase of (g₁) or (g₂).

If $u_1 = wx_i x_j x_i w'$ then $w = w' = \emptyset$ is impossible in virtue of the minimality of U . Suppose $w \neq \emptyset$. Then $u_1 \text{ non } \triangleleft V_6$, $u_2 \triangleleft V_6$ and either $u_2 \in T^*(x^2)$ and we obtain case (f) or $u_2 = x_{m+1} v x_{m+1}$, $v \in T^*(x^2)$, v closed. In the latter case $u_2 \text{ non } \triangleleft V_5$, $u_1 \triangleleft V_5$; hence $w' = \emptyset$. Furthermore, if $v \in T_0(x^2)$, $l(v) = 2m$ then $u_2 \text{ non } \triangleleft V_5^{(m)} = \{y b_n y \mid n > m\}$. Thus, $v \in T_1(x^2)$. Now if $l(w) = 1$ we have case (h₂); if $l(w) > 1$ then $u_1 \text{ non } \triangleleft V_6$, $u_2 \triangleleft V_6$ which implies $v = yv'$ and we obtain (g₂). The case where $w' \neq \emptyset$ can be settled analogously.

Now let $x^2 \text{ non } \triangleleft u_1$, $x^2 \text{ non } \triangleleft u_2$. Put $u_1 \triangleleft V_6$; then $u_1 = x_{m+1} w x_{m+1}$, $w \in X^*$, closed, $l(w) > 1$ by the minimality of U . Now $u_1 \text{ non } \triangleleft V_6$, and thus $u_2 \doteq w u' w'$ where $w, w' \in T(xy x)^0$, $u' \in X^*$. If $w, w' \neq \emptyset$ we have case (c), if $w = \emptyset$ or $w' = \emptyset$ we get a subcase of (g_1) or (g_2) , respectively. This completes the proof of the necessity.

§ 3. If

In proving that the sets given in Theorems 1 and 2 are essentially finite we have to show that infinite sets depending on them are dependent in themselves. For this, we need some theorems which assure the dependence of certain types of infinite sets, and these theorems in their turn are based on some results in the theory of q.o. sets. We are going now to quote these latter ones.

Definition 10. The quasi-ordered set P is a *well quasiordered set* (WQOS) if it satisfies the descending chain condition and does not contain infinite independent subsets (i.e. infinite sets of pairwise incomparable elements).

Next we give some plain facts.

(I) *Let P be a quasi-ordered set. If there exists a mapping γ of P in a WQOS R such that $p\gamma \cong p'\gamma$ implies $p \cong p'$ then P is a WQOS itself.*

Let us mention two important particular cases:

(I₁) *A subset of a WQOS is a WQOS.*

(I₂) *If \prec is a refinement of the q.o. $<$ on P and P is a WQOS under $<$ then so it is under \prec .*

(II) *The union of a finite number of WQOS's is a WQOS.*

Now let P be a q.o. set. Define a q.o. on the set \bar{P} of all finite sequences of elements of P by

$$\pi = (p_1, \dots, p_n) \cong (p'_1, \dots, p'_m) = \pi'$$

iff there exists a subsequence $(p'_{i_1}, \dots, p'_{i_n})$, $1 \cong i_1 < \dots < i_n \cong m$, of π' such that $p_j \cong p'_{i_j}$. The following proposition is a consequence of a theorem of G. HIGMAN [2].

(III) *If P is a WQOS then so is \bar{P} .*

We prefer to give here a self-contained proof. First of all, the classical theorem of RAMSEY implies:

(IV) *The direct product of a finite number of WQOS's is a WQOS.*

It is routine to check the validity of the descending chain condition in \bar{P} . Now suppose P_1 is an infinite independent subset of \bar{P} and let $\pi = (p_1, \dots, p_n)$ be an element of minimal length in P_1 . Suppose $\pi' = (p_1, \dots, p_{k-1})$ is the maximal segment of π such that the subset $R = \{\varrho \in P_1 \mid \pi' < \varrho\}$ of P_1 is infinite; obviously, $0 < k \cong n$. Choose

a subsequence $q' = (r_{q(1)}, \dots, r_{q(k-1)})$ in each $q \in R$ such that $p_i \cong r_{q(i)}$, each $q(i)$ as small as possible. By (IV) and the theorem of RAMSEY there exists an infinite subset R' of R such that the set $R'' = \{q' \mid q \in R'\}$ is totally quasi ordered (i.e. any two elements of R'' are comparable).

Now we "break up" the elements of R' : for each $q \in R'$ put $q^{(i)} = \{r_{q(i-1)+1}, \dots, r_{q(i)-1}\}$ where $q(0) = 0$, $q(k) = l(q) + 1$ and $q(i)$ is defined as above if $1 \leq i \leq k-1$. It can be seen easily that every component of $q^{(i)}$ ($i = 1, \dots, k$) is either strictly less than p_i or incomparable with it. At least one of the sets $R^{(i)} = \{q^{(i)} \mid q \in R'\}$ must contain an infinite independent subset; indeed, in the opposite case all they are WQOS's and, by (IV), so is the direct product $Q = R'' \oplus R^{(1)} \oplus \dots \oplus R^{(k)}$. However then the mapping $\gamma: R' \rightarrow Q$ ($q\gamma = (q', q^{(1)}, \dots, q^{(k)})$) satisfies the conditions in (I) and therefore R' is a WQOS which is a contradiction.

Suppose $R^{(i)}$ contains an infinite independent subset P_2 . Repeating the above construction, we find a component $p_{i_1 i_2}$ of a vector of minimal length in P_2 such that there exists an infinite independent set P_3 consisting of vectors with components strictly less than or incomparable with $p_{i_1 i_2}$. Thus we obtain an infinite series $p_{i_1}, p_{i_1 i_2}, \dots$ of elements of P having the property that every member of it is either strictly less than or incomparable with each of the preceding ones which is impossible.

In applying these facts to word sets we shall use besides \triangleleft three further relations:

Definition 1₁. $u \triangleleft_1 v$ iff there exists an endomorphism $\varphi: F \rightarrow F$ such that $u\varphi$ is an end part of v ($v = v' \cdot u\varphi$).

Definition 1_r. $u \triangleleft_r v$ iff there exists an endomorphism $\varphi: F \rightarrow F$ such that $u\varphi$ is a beginning part of v ($v = u\varphi \cdot v'$).

Definition 1_q. $u \triangleleft_q v$ iff there exists an endomorphism $\varphi: F \rightarrow F$ such that $v = u\varphi$.

For a set of words V the q.o. set $\{V, \triangleleft\}$ will be simply denoted by V ; furthermore, we shall write

$$V_l = \{V, \triangleleft_l\}, \quad V_r = \{V, \triangleleft_r\}, \quad V_q = \{V, \triangleleft_q\}.$$

Obviously, all these sets satisfy d.c.c. By (I₂), if V_q is a WQOS then so is V_l and V_r , and if either of these latter ones is a WQOS then so is V .

We introduce an operation \circ on sets of words as follows:

$$U \circ V = \{w \mid w \equiv uv, u \in U, v \in V\}.$$

It holds

Theorem 3. Let $U, U', V_r, V_l', W_q, W_q'$ be WQOS's. Then $V' \circ V, (V' \circ W)_l, (W' \circ V)_r$, and $(W' \circ W)_q$ are WQOS's. Moreover, if either the last letter of every element of U and V or the first letter of every element of U' and V' is closed then $U \circ U', (U \circ W)_l, (W \circ U)',$ and $(V \circ V')_q$ are WQOS's, too.

Proof. Let us prove the assertion concerning $U \circ U'$; in the other cases the proof runs analogously. Suppose for example that $U = \{u_i | u_i \doteq u_i^* y_i, i = 1, 2, \dots\}$, $U' = \{u'_j | j = 1, 2, \dots\}$. Then $U \circ U' = \{w_{ij} | w_{ij} \doteq u_i u'_j\}$. Now the direct product $U \times U'$ is a WQOS by (IV). Consider the mapping $\gamma: U \circ U' \rightarrow U \times U'$ defined by $w_{ij} \gamma = (u_i, u'_j)$. Now let $(u_i, u'_j) \cong (u_k, u'_p)$ i.e. $u_i \triangleleft u_k, u'_j \triangleleft u'_p$ and, say, $u_k = s(u_i \varphi) t = s(u_i^* \varphi)(y_i \varphi) t$, $u'_p = s'(u'_j \psi) t'$. As u_i^* and u'_j are relatively prime and none of them contains y_i , there exists an endomorphism $\chi: F \rightarrow F$ such that $u_i^* \chi = u_i^* \varphi, u'_j \chi = u'_j \psi$ and $y_i \chi = y_i \varphi \cdot t s'$. Thus $w_{ki} = u_k u'_i = s w_{ij} \chi t'$, i.e. $w_{ij} \triangleleft w_{ki}$. Hence the assertion follows by (I).

Theorem 4. Let T be a type. Then $T^*, T_1^*, T_r^*, T_{1q}$ are WQOS's.

Proof. First remark that X_q^* and T_{0r} are WQOS's. Hence, by Theorem 3 (case $(V \circ V')_q$) and (II), the posets $T'_q = (T_0 \circ X^*)_q$ and $T''_q = (T' \cup X^*)_q$ are WQOS's. Using the notation $T'^* = \{w | w = \prod_{i=1}^n t_i, t_i \in T'\}$, we have $T_1 = X^* \cup T'^* \cup (X^* \circ T'^*) \cup (T'^* \cup (X^* \circ T'^*)) \circ T_0$. Thus, it suffices to show that T_1^* is a WQOS since then by Theorem 3 $(T' \circ T'^*)_q$ and hence $(T' \cup (T' \circ T'^*))_q = T''_q$ are also WQOS's. Using again Theorem 3, (I₂) and (II) we conclude, furthermore, that T_{1q} and $T_r^* = (T_0 \cup T_1)_r$ are WQOS's, too. The rest follows by duality and by (I₂), respectively.

Since T'_q is a WQOS, the same follows from (III) for the set of finite sequences $\overline{T'_q}$. Consider the mapping $\gamma: T_1'^* \rightarrow \overline{T'_q}$ defined by $w \gamma \doteq \left(\prod_{i=1}^n t_i \right) \gamma = (t_1, \dots, t_n) (t_i \in T')$.

We have to show that γ satisfies the condition of (I). Put $w' = \prod_{j=1}^m t'_j \in T'^*$ and suppose $(t_1, \dots, t_n) \cong (t'_1, \dots, t'_m)$ in $\overline{T'_q}$, i.e. $t_i \triangleleft_q t'_{j_i}$ for some $1 \leq j_1 < \dots < j_n \leq m$. In other words, there exist endomorphisms $\varphi_i: F \rightarrow F$ ($i = 1, \dots, n$) such that $t_i \varphi_i = t'_{j_i}$. Denote the last letter of t_i by y_i ; then $t_i \doteq t_i^* y_i$. Since t_1^*, \dots, t_n^* are pairwise relatively prime and they do not contain y_1, \dots, y_n , there exists an endomorphism $\varphi: F \rightarrow F$ such that $t_i^* \varphi = t_i^* \varphi_i$ for $i = 1, \dots, n$; $y_i \varphi = y_i \varphi_i \cdot \prod_{k=j_i+1}^{j_{i+1}-1} t'_k$ for $i = 1, \dots, n-1$ and $y_n \varphi = y_n \varphi_n \cdot \prod_{k=j_n+1}^m t'_k$. Hence $t_i \varphi = \prod_{k=j_i}^{j_{i+1}-1} t'_k$ for $i < n, t_n \varphi = \prod_{k=j_n}^m t'_k$ and $w' = \left(\prod_{k=1}^{j_1-1} t'_k \right) \times \dots \times (w \varphi)$. The theorem is proved.

Theorem 5. Let $V = \{v_i | i = 1, 2, \dots\}$ be a set of words. Suppose there exist natural numbers k, l, n and n types $T^{(1)}, \dots, T^{(n)}$ such that every v_i has a decomposition $v_i = v_{i0} \cdot \prod_{j=1}^n u_{ij} v_{ij}$ where

- a) $l(u_{ij}) \leq l$,
- b) $v_{ij} \in T^{(j)*0}$,
- c) v_{ij} is closed if non-empty,
- d) if $l(v_{ij}) > k$ for some $j \neq 0, n$ then $v_{ij} \in T_1^{(j)}$.

Then V is dependent. Moreover, if $v_{i_0} \neq \emptyset$ ($v_{i_n} \neq \emptyset$) for every i then V_r (V_l) is dependent as well.

Proof. We shall prove that V contains an infinite WQOS. Indeed, there is only a finite number of sequences u_1, \dots, u_n of words of bounded length. Consequently, V has an infinite subset V' such that to every $v_i \in V'$ the same sequence u_1, \dots, u_n corresponds. Now $V' = V_0^* \cup V_0^{**}$ where $V_0^* = \{v_i \mid v_i \in V', v_{i_0} \neq \emptyset\}$, $V_0^{**} = V' \setminus V_0^*$. Put $V'_0 = V_0^*$ if V_0^* is infinite and $V'_0 = V_0^{**}$ in the opposite case. Construct in a similar way consecutively $V'_0 \supseteq V'_1 \supseteq \dots \supseteq V'_n$ with either $V'_r = V_r^* = \{v_i \mid v_i \in V'_{r-1}, v_{i_r} \neq \emptyset\}$ or $V'_r = V'_{r-1} \setminus V_r^*$. For sake of simplicity suppose $u_j \neq \emptyset$ and, for $j \neq 0, j \neq n$, $V'_j = V_j^*$ (in the opposite case we possibly had to change the parameters k, l, n and the decomposition of v_i). Put $T_0^{(j)}(k) = \{w \mid w \in T_0^{(j)}, l(w) \leq k\}$, $U_j = T_1^{(j)} \cup T_0^{(j)}(k)$ for $1 \leq j \leq n-1$, $U_0 = T^{(0)*}$ or $\{\emptyset\}$ according to $V'_0 = V_0^*$ or V_0^{**} and $U_n = T^{(n)*}$ or $\{\emptyset\}$ according to $V'_n = V_n^*$ or V_n^{**} . Then $v_{ij} \in U_j$ for $v_i \in V'_n$, and U_{j_q} ($1 \leq j \leq n-1$) as well as U_{0l}, U_{nr} (if different from $\{\emptyset\}$) are WQOS's. This implies that $A = U_{0l} \times \left(\prod_{j=1}^{n-1} U_{j_q} \right) \times U_{nr}$ is also a WQOS. Define $\gamma: V'_n \rightarrow A$ by $v_i \gamma = \left(v_{i_0} \cdot \prod_{j=1}^n u_j v_{i_j} \right) \gamma = (v_{i_0}, \dots, v_{i_n})$. Suppose $v_h \in V'_n$ and $(v_{i_0}, \dots, v_{i_n}) \leq (v_{h_0}, \dots, v_{h_n})$ in A , i.e. $v_{h_0} = w \cdot (v_{i_0} \varphi_0)$, $v_{h_n} = (v_{i_n} \varphi_n) \cdot w'$ (if non-empty) and $v_{h_j} = v_{i_j} \varphi_j$ ($j = 1, \dots, n-1$) with some suitable endomorphisms $\varphi_0, \dots, \varphi_n$ of F . In consequence of c), there exists $\varphi: F \rightarrow F$ such that $u_j \varphi = u_j$, $v_{i_j} \varphi = v_{i_j} \varphi_j$ so that $v_h = w(v_i \varphi)w'$. Consequently, V_n is an infinite WQOS and therefore contains comparable elements which completes the proof.

Now we are in position to prove the second parts of Theorems 1 and 2.

Sufficiency (Theorem 1). In consequence of Proposition 6, every infinite set depends on $u \in X^*$. If $u \in T(xy^2x)$ and u non $\triangleleft V$ then, by Proposition 7, $V \subseteq T^*(x^2)$ and it is dependent in itself by Theorem 4. This completes the proof of Theorem 1.

Sufficiency (Theorem 2). We shall show that U is an EFS in cases (c), (e₁), (f), (g₁), (h₁) and in the subcase of (d) where $v, v' \in T(x^2)$. Cases (e₂), (g₂), (h₂) follow then by duality, (d') and the rest of (d) from Proposition 4, and (a), (b) from Proposition 3. In all cases the proof consists in finding the general form of words which are independent from U and in a subsequent application of Theorem 5 to infinite sets consisting of such words. Thus, put U non $\triangleleft q$. If $q = q_1 x q_2$ we shall say that x occurs in q later (earlier) again if $q_2 = q' x q''$, $q' \neq \emptyset$ ($q_1 = q' x q''$, $q'' \neq \emptyset$).

Case (c). Let $q = q' t q''$ where $q', q'' \in T^*(x^2)^0$ and either $t = \emptyset$ or the first letter of t occurs later again, the last one earlier again in q . We have $l(t) < 2l(w) + l(w') + 2$. Indeed, the number of letters between the first and last occurrence of the extreme letters of t cannot exceed $l(w) - 1$, and the number of those between the last occurrence of the first letter of t and the first occurrence of the last one must be less than $l(w')$. Put k arbitrary, $l = 2l(w) + l(w') + 1$, $n = 1$, $T^{(0)} = T^{(1)} = T(x^2)$ and apply Theorem 5.

Case (d), $v, v' \in T(x^2)$. Let $q \doteq q'tq''$ where $q'q'' \in T^*(xyx)^0$, $t = x_i^2 t' x_j^2$ or $t \in T(x^2)^0$ (as $xyzx \text{ non } \triangleleft q$, such a decomposition exists by Proposition 8). Then $l(t') < l(w)$ and we can apply Theorem 5 with k arbitrary, $l = l(w) + 3$, $n = 1$, $T^{(0)} = T^{(1)} = T(xyx)$.

Case (e₁). Since $xyzx \text{ non } \triangleleft q$, there exists a decomposition $q \doteq q'tq''$ where $q' \in T^*(xyx)^0$, $q'' \in T^*(x^2)^0$ and $t = x_i^2 t' x x_j x_k x_j$ (or $t = \emptyset$). Then $l(t') < l(w)$. Put k arbitrary, $l = l(w) + 4$, $n = 1$, $T^{(0)} = T(xyx)$, $T^{(1)} = T(x^2)$ and Theorem 5 yields the dependence.

Case (f). If $q = q_1 x_j q_2 x_j q_3$ with $q_2 \neq \emptyset$ then either $l(q_1) < l(w)$ or $l(q_3) < l(w')$. Hence $q = tq't'$ where $l(t) \leq l(w)$, $l(t') \leq l(w')$ and $q' \in T^*(x^2)^0$; moreover, if x_{i_1}, \dots, x_{i_s} are the letters which occur in q' and in either t or t' , too, then $q = tv_0 \left(\prod_{j=1}^s x_{i_j}^{\varepsilon_j} v_j \right) t'$ where $\varepsilon_j = 1$ or 2 , $v_j \in T^*(x^2)^0$, v_j closed if non-empty and if $l(v_j) \geq 2m(v)$ ($m(v)$ denotes the number of different letters in v) then $v_j \in T_1(x^2)$ (in the opposite case we had $v \triangleleft q$). Thus, we can apply Theorem 5 putting $k = 2m(v)$, $l = 2(l(w) + l(w'))$, some $n \leq l(w) + l(w') + 2$ and $T^{(0)} = \dots = T^{(n)} = T(x^2)$.

Case (g₁). Here $q = v_0 \left(\prod_{j=1}^s x_{i_j}^{\varepsilon_j} v_j \right) t$ where $l(t) = l(w)$, x_{i_1}, \dots, x_{i_s} are different letters occurring also in t , $\varepsilon_j = 1$ or 2 , $v_j \in T^*(x^2)^0$, v_j closed if non-empty and $v_j \in T_1(x^2)$ if $l(v_j) > 2m(v)$, $j \neq 0$ (in the opposite case $xvyx \triangleleft q$). One can apply Theorem 5 with $k = 2m(v) + 1$, $l = 2l(w)$, some $n \leq l(w) + 1$ and $T^{(0)} = \dots = T^{(n)} = T(x^2)$.

Case (h₁). Now either $q \in T^*(x^2)$ or $q = v_0 x_i v_1 x_i$ where $\varepsilon = 1$ or 2 , $v_0, v_1 \in T^*(x^2)^0$, closed, and $v_1 \in T_1(x^2)$ if $l(v_1) > 2m(v)$. Again, put $k = 2m(v)$, $l = 2$, $n = 2$, $T^{(0)} = T^{(1)} = T^{(2)} = T(x^2)$ (as a matter of fact, $T^{(2)}$ is irrelevant since $U_2 = \emptyset$).

Theorem 2 is proved.

Corollary. If the identity $u = u'$ defines a hereditarily finitely based variety then the pair of types $T(u)$, $T(u')$ is one of the pairs (a)—(h₂).

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