## Schedule on parallel machines in the case of individual machine-set

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There is a less researched area of the parallel machines scheduling, when there is an $M_{j}$ machine-set for each job $j$, where it can be scheduled. We would like to minimize the latest job's finishing time, that is the makespan, $C_{\max }$ so, that each job can be processed only by one machine at a time and one machine can process at most one job at a time. This problem is $\mathcal{N}$ - - complete, since its special case, the $P \| C_{\text {max }}$ problem, (when each job can be processed by each machine), is known to be $\mathcal{N P}$-complete [1]. If we have restrictions for the $M_{j}$ machine-sets, the well-known list-schedule gives a nearly optimal solution. We give a better approximation algorithm for the optional problem.
M. Pinedo [2] has studied that special case of the problem, when the job's processing times are 1, and he has showed that,if the $M_{j}$ sets are laminals, then an easy list-schedule gives an optimal solution. We show for the general case, when the $M_{j}$ sets are optional ( $P\left|M_{j}, p_{j}=1\right| C_{\text {max }}$ ), that it is a network-flow-problem. Moreover, from this, we have an algorythm for that special case, when we allow the preemtion for the jobs $\left(P\left|M_{j}, p m t n\right| C_{m a x}\right)$. From this algorithm we have received a minimax formula for the optimal makespan too.
We study the problem when we allow a special preemption, called $p m t n^{*}$, when a job can be split, but it should be processed immediately by another machine. This special preemption gives a better solution, that we show in an example, that the optimal value of the problem $P\left|M_{j}, p m t n^{*}\right| C_{\max }$ is less than the equivalent $P\left|M_{j}\right| C_{\max }$ problems optimal value. But we can observe, that if $M_{j}=M$ for each job, than the special preemption gives no better optimal value, so in this case the problem is equivalent with the $P \| C_{\text {max }}$ problem, which is $\mathcal{N} \mathcal{P}$-complete. We have a 2-approximate algorythm for this problem too.

## References

[1] M.R.Garey, D.S. Johnson [1978]: Strong NP- completeness results: motivaton, examples and implications, Journal of the Association for Computing Machinery 25, 499-508
[2] M.Pinedo [2002]: Scheduling Theory, Algorithms, and Systems, Second Edition, Prentice Hall

